

Bosonic Dirty Paper Coding

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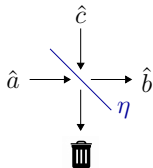
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- The bosonic (Gaussian) channel is a simple quantum-mechanical model for optical communication over free space or optical fibers

Bosonic Model

For a single-mode lossy bosonic channel, the channel input is an electromagnetic field mode with annihilation operator \hat{a} , and the output is

$$\hat{b} = \sqrt{\eta} \hat{a} + \sqrt{1 - \eta} \hat{e}$$



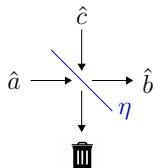
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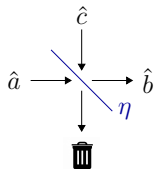
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- the noise mode \hat{e} is in a thermal Gaussian state
- the transmissivity $\eta \in [0, 1]$ depends on the absorption length of the optical fiber

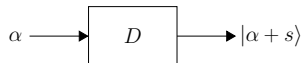


- The transmitter employs a coherent state protocol.
- A coherent state $|\alpha\rangle$ corresponds to an oscillation of the electromagnetic field,

$$|\alpha\rangle = D(\alpha)|0\rangle$$

$$D(\alpha) \equiv \exp(\alpha \hat{a}^\dagger - \alpha^* \hat{a})$$

We consider the single-mode lossy bosonic channel with a coherent-state protocol and non-ideal modulation:



■ Applications

- classical interference in the transmission equipment
- watermarking with a quantum embedding

- Homodyne and heterodyne detection
- Joint detection
 - DPC lower bound
 - MMSE coefficient is sub-optimal

In the Gel'fand-Pinsker model,

- the channel depends on a **random parameter**, $S \sim p_S$, while the transmitter has channel side information (CSI).
- Applications:
 - cognitive radio in wireless systems
 - memory storage
 - digital watermarking

Theorem (GP, 1980; Heegard and El Gamal, 1983)

The capacity of a random-parameter classical channel $W_{Y|X,S}$ with CSI at the transmitter is given by

$$C(W) = \max_{P_{U,X|S}} (I(U; Y) - I(U; S))$$

where $U \rightarrow (X, S) \rightarrow Y$ form a Markov chain.

Gaussian Channel with Additive Interference

Consider

$$Y = X + S + Z$$

with Gaussian noise $Z \sim \mathcal{N}_{\mathbb{R}}(0, \sigma^2)$ and interference S .

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The capacity of the Gaussian channel with additive interference and CSI at the transmitter is the same as if there is no interference, i.e.

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- It is not obvious. The trivial strategy is to send $X = U - S$ such that $U \perp S$, resulting in an interference-free output, $Y = U + Z$. However, this strategy would waste transmission power.

Dirty Paper Coding (DPC) Strategy

Set

$$U = X + tS,$$

where $X \sim \mathcal{N}_{\mathbb{R}}(0, P)$, such that X is statistically **independent** of S .

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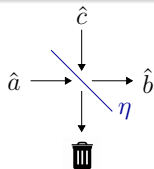
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- Explicit lattice codes were proposed e.g. by [Erez and ten Brink, 2005].

Lossy Bosonic Channel

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$$\hat{b} = \sqrt{\eta} \hat{a} + \sqrt{1 - \eta} \hat{e}$$



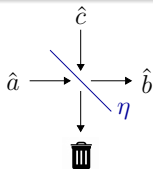
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where

- the noise mode \hat{e} is in a thermal state, $\tau(N_E) \equiv \int_{\mathbb{C}} d^2\alpha \frac{e^{-|\alpha|^2/2N_E}}{\pi N_E} |\alpha\rangle\langle\alpha|$



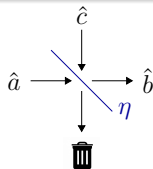
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- $0 \leq \eta \leq 1$ is the transmissivity, which depends on the absorption length of the optical fiber



Coding with CSI

- Alice chooses $m \in \{1, 2, \dots, M\}$

Encoding: $m \mapsto (\alpha_i(m, s_1, \dots, s_n))_{i=1}^n, |\alpha_i| \leq N_A.$

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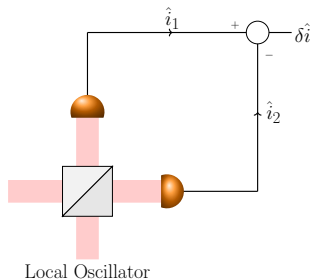
The coding rate is defined as $R = \frac{\log(M)}{n}$ [bits per transmission], and the maximal probability of error is denoted by $P_e^{(n)} = \max_m \Pr(\hat{M} \neq m | m)$. A rate $R > 0$ is called achievable if there exists a sequence of $(\lceil 2^{nR} \rceil, n)$ codes such that $P_e^{(n)} \rightarrow 0$ as $n \rightarrow \infty$. The **operational capacity** $C(\mathcal{E})$ is defined as the supremum of achievable rates.

This can be viewed as a watermarking model with a quantum embedding.

- Given a classical host data sequence s_1, \dots, s_n , Alice encodes an authentication message m into a watermark $(\alpha_i(m, s_1, \dots, s_n))_{i=1}^n$.
- Next, she performs a quantum embedding of the watermark; she prepares a *watermarked state* $|\alpha_i + s_i\rangle$ ("stegotext") and transmits it to Bob through the optical fiber.
- The capacity of the random-parameter bosonic channel represents the optimal rate at which the authenticator, Bob, can recover the messages with high fidelity.

Homodyne Measurement

A homodyne measurement of a quadrature observable is implemented by combining the target quantum mode with an intense local oscillator at a 50:50 beam splitter, and measuring the photocurrent difference of the outgoing modes using two photodetectors.



Homodyne Measurement (Cont.)

When homodyne detection is used with a coherent-state protocol, the resulting channel \mathcal{E}_{hom} is the classical Gaussian channel

$$Y = \sqrt{\eta}(\alpha + S) + Z_{\text{hom}}$$

with a real-valued $S \sim \mathcal{N}_{\mathbb{R}}(0, N_S)$ and noise $Z_{\text{hom}} \sim \mathcal{N}_{\mathbb{R}}\left(0, \frac{1}{4} [2(1 - \eta)N_E + 1]\right)$

Results: Homodyne Measurement

Using the DPC scheme, $\alpha \sim \mathcal{N}_{\mathbb{R}}(0, N_A)$ and

$$U = \alpha + t_0 S$$

such that α and S are uncorrelated, with $t_0 = \frac{\eta N_A}{\eta N_A + \text{var}(Z_{\text{hom}})}$.

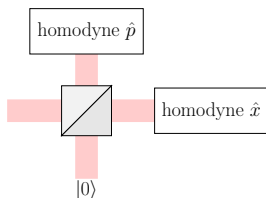
The effect of the interference is thus removed, and the capacity is given by

$$C(\mathcal{E}_{\text{hom}}) = \frac{1}{2} \log_2 \left(1 + \frac{4\eta N_A}{2(1-\eta)N_E + 1} \right)$$

as without interference.

Heterodyne Measurement

In heterodyne detection, two quadratures are measured by combining the measured mode with a vacuum mode into a 50:50 beam splitter, and homodyning the quadratures of the outcome modes.



Heterodyne Measurement (Cont.)

Heterodyne detection is described by a random-parameter channel \mathcal{E}_{het} with complex-valued Gaussian noise,

$$Y = \sqrt{\eta}(\alpha + S) + Z_{\text{het}}$$

with a complex-valued circularly-symmetric Gaussian random parameter $S \sim \mathcal{N}_{\mathbb{C}}(0, \frac{1}{2}N_S)$ and noise $Z_{\text{het}} \sim \mathcal{N}_{\mathbb{C}}(0, \frac{1}{2}[(1 - \eta)N_E + 1])$.

Results: Heterodyne Measurement

Using the DPC scheme, $\alpha \sim \mathcal{N}_{\mathbb{C}}(0, \frac{1}{2}N_A)$ and

$$U = \alpha + t_1 S$$

such that α and S are uncorrelated, with $t_1 = \frac{\eta N_A}{\eta N_A + \text{var}(Z_{\text{het}})}$.

The capacity is given by

$$C(\mathcal{E}_{\text{het}}) = \log \left(1 + \frac{\eta N_A}{(1 - \eta) N_E + 1} \right)$$

the quantum counterpart of the classical channel with additive white Gaussian noise (AWGN)

Results: Joint Detection

For joint detection, the channel does not have a classical description.

Applying the previous result in [P., 2020] for a quantum channel with random parameters, and using the DPC strategy, we obtain the lower bound $C(\mathcal{E}_{\text{joint}}) \geq R_{\text{DPC}}(t)$,

$$R_{\text{DPC}}(t) \equiv I(\gamma; B) - I(\gamma; S) \Big|_{\gamma=\alpha+tS}$$

DPC Lower Bound

$$R_{\text{DPC}}(t) = g(\eta(N_A + N_S) + (1 - \eta)N_E) \\ - g\left(\frac{\eta(1 - t)^2 N_A N_S}{N_A + t^2 N_S} + (1 - \eta)N_E\right) - \log_2\left(\frac{N_A + t^2 N_S}{N_A}\right)$$

where

$$g(N) = \begin{cases} (N + 1) \log_2(N + 1) - N \log_2(N) & N > 0 \\ 0 & N = 0. \end{cases}$$

Pure-Loss Bosonic Channel

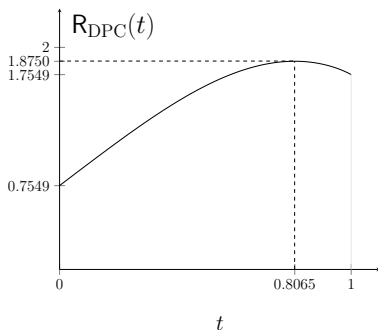
For $N_E \rightarrow 0$,

$$R_{\text{DPC}}(t) = g(\eta(N_A + N_S)) - g\left(\frac{\eta(1-t)^2 N_A N_S}{N_A + t^2 N_S}\right) - \log\left(\frac{N_A + t^2 N_S}{N_A}\right)$$

Results: Joint Detection (Cont.)

For example, suppose that $N_A = N_S = 2$ and $\eta = \frac{1}{2}$. Then,

$$R_{\text{DPC}}(t) = g(2) - g\left(\frac{(1-t)^2}{1+t^2}\right) - \log(1+t^2)$$



Achievable Rates

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- Optimal DPC coefficient is $t_{\text{max}} = 0.8065 \Rightarrow R_{\text{DPC}}(t_{\text{max}}) = 1.8750$
- joint-detection capacity without interference: $g(1) = 2$.

Thank you